

APPENDIX A

SLOOP SYNTAX QUICK REFERENCE

A.1 Notational conventions

The SLOOP syntax is given below in BNF. The significance of the typeface and symbols used in the BNF description is as follows:

Plain or boldface type	Terminal symbols
Italics	Non-terminal symbols
Braces	Grouping
Square brackets	Enclosed syntactical unit is optional
Asterisk	Zero or more occurrences of the syntactical unit
Plus	One or more occurrences of the syntactical unit
Vertical bar	Separates options
Single quotes	Enclose literals.

All names, such as *program-name*, *class-name*, *instance-name* and *category-name*, are strings that may contain letters, digits and the underscore character. They all start with a letter. If the plural form is used (for example as in *instance-variable-names*), then one or more name(s) may be present, each separated by a white-space character. Three consecutive colons separate a *class-name* from a *package-name* when it is necessary to qualify the *class-name* by a *package-name*.

A.2 The SLOOP program structure

A *SLOOP-program* has the following structure :

<i>SLOOP-program</i>	→	program <i>program-name</i> <i>activation-section</i> { <i>package-description</i> }+ end-program
<i>activation-section</i>	→	sequential <i>statement-list</i> end-sequential parallel <i>statement-list</i> end-parallel

package-description → **package** *package-name*
 {*package-description*}*
 {*class-description*}*
 {*partial-class-description*}*
end-package

A.3 Complete and partial class descriptions

Syntax of a *class-description*:

class-description →
class *class-name*
superclass *superclass-name* [**from** *repository-name*]
 [**class variable names** [*class-variable-names*]]
 [**instance variable names** [*instance-variable-names*]]
 [**class macros** [*macros-section*]]
class properties [*properties-section*]
methods-section

Syntax of a *partial-class-description*:

partial-class-description →
class *class-name* **from** *repository-name*
partial-class-methods-section

The **class properties** keywords are mandatory in order to indicate to the designer that all the relevant properties should always be listed. If this section had been optional, then it would have been unclear to the person reusing the class whether the original designer had merely chosen not to list the properties or whether there had been no relevant properties to list.

A.4 The *macros-section*

Syntax of a *macros-section*:

macros-section → *macro-list*
macro-list → *macro-definition*
 {[] *macro-definition*}*
macro-definition → *macro-variable* ≡ *macro-expression*
macro-variable → *variable-name*
macro-expression → *simple-macro-expression* |
conditional-macro-expression
simple-macro-expression → *message-expression* | *variable-name* | *literal*
conditional-macro-expression → *simple-macro-expression*
 if *boolean-expression*
 {~ *simple-macro-expression*
 if *boolean-expression*}*

A *message-expression* is a Smalltalk-style message expression and a *boolean-expression* is a Smalltalk-style message expression that returns true or false. A Smalltalk-style message expression consists of a *receiver*, a *selector* and zero or more *arguments*. If there are no arguments, the message is called a *unary message*. For example, `bufferedElements size` is

a unary message expression. A *binary message* has a single argument following a selector consisting of one or two non-alphanumeric characters, the second of which may not be a minus sign. The message expression `a + b` is an example of a binary message expression. The third type of message is a *keyword message*. The selector consists of one or more keywords, each with its associated argument. A keyword consists of an identifier followed by a colon. For example, `bufferedElements addLast: newElement` is a keyword message expression.

A *literal* is a Smalltalk-style literal which may be a number, a symbol constant, a character constant, a string or an array constant.

Note that, as in Smalltalk-80, message expressions may be nested. The receiver of a message expression may itself be a message expression. Similarly, the argument(s) of a keyword message may also be message expression(s).

A message expression may also contain a block. The reader is referred to [GoRo89] for a detailed discussion of a Smalltalk-80 block. For the purposes of its application in SLOOP, the following description suffices.

A block represents a deferred sequence of actions. It consists of a sequence of expressions separated by periods and delimited by square brackets. The actions represented by a block are not necessarily executed when the block expression is encountered. For example, a block expression is used as the argument of the Smalltalk-80 `and:` keyword message. This message represents the logical "and" operation. If the first operand (i.e. the receiver) evaluates to true, the second operand (the argument of `and:`) is evaluated and its value is returned as result. However, if the first operand evaluates to false, the second operand is not evaluated. This is indicated syntactically via the fact that the argument of the `and:` message is a block expression.

A block is also used when the receiver of a message is a collection and the actions represented by the block need to be applied to each element of the receiver. In that case each element of the receiver is passed to the block as an argument. This is indicated syntactically by the presence of an identifier preceded by a colon at the beginning of the block. This identifier is separated from the rest of the contents of the block by a vertical bar.

For example, the Smalltalk-80 library `select:` and `detect:` messages are used frequently in the `CallCentreSimulation` example. The `select:` message evaluates the block received as argument of the message for each of the receiver's elements (the receiver is a collection object). It returns a collection that contains only those elements of the receiver for which the block evaluates to true. The message expression below returns the collection representing all employees earning a salary greater than \$20 000:

```
employees select: [:each | each salary > 20000]
```

The following message expression returns the object representing the first employee found earning a salary greater than \$20 000 (if no such employee is found, then `nil` is returned):

```
employees detect: [:each | each salary > 20000] ifNone: [nil]
```

For a formal description of the Smalltalk-80 syntax, the reader is referred to [GoRo89].

A.5 The *properties-section*

Each class, as well as each method within a class may contain a *properties-section*. The syntax is as follows:

properties-section → *property-list*
properties-list → *property*
 {[] *property*}*

A *property* may be of the form:

p unless q,
stable *p*,
invariant *p*,
p ensures q,
p leads-to q,
p until q,
p detects q,
p precedes q or
p results-in q,
 where *p* and *q* are **first-order** predicates.

Since **first order predicate logic** is used, universal and existential quantification is allowed in the logical relations. The keywords **forall** and **exists** may be used as alternatives to the \forall (universal quantification) and \exists (existential quantification) symbols respectively. Instead of using a colon to denote the domain of a quantification, the reserved word **where** is used. This is to avoid confusion with the colon used in Smalltalk keyword expressions.

If a variable-list is used in a quantification, the variables are separated by a comma preceded by a backslash. This ensures that the comma cannot be mistaken for the Smalltalk concatenation symbol. If a '<' symbol is followed immediately by a quantification symbol, it denotes the start of a quantification construct. If a '>' symbol appears as the first non-white-space character on a line, it denotes the end of a quantification construct. The quantification constructs have the same parsing precedence as parentheses.

Smalltalk-style message expressions may be used in the predicates. However, these message expressions may not have any side-effects, i.e. they may not change the state of any object.

The \wedge (logical and), \vee (logical or) and \neg (negation) operators are defined in addition to the Smalltalk **&** (logical and), **|** (logical or) and **not** (negation) operators. The additional logical operators serve a readability purpose only. One difference between the additional logical operators and the Smalltalk ones is in the parsing precedence. The Smalltalk unary, binary and keyword expressions are evaluated in that order. The additional logical operators are evaluated after the unary, binary and keyword expressions have been evaluated.

Further conventions about the priorities of logical relations are given next (those on the same line have equal priority and the lines represent the priorities from high to low):

\neg

$=, \neq$

\wedge, \vee

\Rightarrow

\equiv

unless, ensures, leads-to, stable, invariant, detects, until, precedes

Properties are universally quantified over all the free variables occurring in them.

A.6 The *methods-section*

The syntax of the *methods-section* of a class is as follows:

<i>methods-section</i>	→	[class methods { <i>methods-implementation</i> }+] [instance methods { <i>methods-implementation</i> }+]
<i>methods-implementation</i>	→	category <i>category-name</i> { <i>method-description</i> }+
<i>method-description</i>	→	<i>sequential-method</i> <i>parallel-method</i>
<i>sequential-method</i>	→	message pattern <i>Smalltalk-message-pattern</i> [method macros <i>macros-section</i>] method properties <i>properties-section</i> sequential <i>statement-list</i> end-sequential
<i>parallel-method</i>	→	message pattern p_ <i>Smalltalk-message-pattern</i> [method macros <i>macros-section</i>] method properties <i>properties-section</i> parallel <i>statement-list</i> end-parallel
<i>partial-class-methods-section</i>	→	[class methods { <i>selector</i> }+] [instance methods { <i>selector</i> }+]
<i>selector</i>	→	<i>Smalltalk-selector</i> p_ <i>Smalltalk-selector</i>

A *Smalltalk-message-pattern* has the usual Smalltalk syntax, i.e. it comprises the message selector with the associated pseudo-variables to represent the arguments if there are any. A *Smalltalk-selector* has the usual syntax of a selector in a Smalltalk program.

Similarly to the properties of the class, all the relevant properties of the method should be listed if there are any. If this section had been optional, then it would have been unclear to the person reusing the class whether the original designer had merely chosen not to list the properties or whether there had been no relevant properties to list.

A.7 SLOOP statements

The SLOOP *statement-list* is defined as follows:

<i>statement-list</i>	→	<i>statement</i> ¹ {[] <i>statement</i> }*
<i>statement</i>	→	<i>simple-statement</i> <i>quantified-statement-list</i>
<i>quantified-statement-list</i>	→	<[] <i>quantification statement-list</i> >

¹ This implies that, as in UNITY, a *statement-list* cannot be empty.



<i>quantification</i>	→	<i>variable-list</i> where <i>boolean-expr</i> ::
<i>variable-list</i>	→	<i>variable</i> {\, <i>variable</i> }*
<i>simple-statement</i>	→	<i>statement-component</i> { <i>statement-component</i> }*
<i>statement-component</i>	→	<i>enumerated-component</i> <i>quantified-component</i>
<i>enumerated-component</i>	→	<i>component-part</i> <i>conditional-component-part-list</i>
<i>quantified-component</i>	→	< <i>quantification</i> <i>simple-statement</i> >
<i>component-part</i>	→	{[[^]] <i>variable</i> := <i>simple-expr</i> } ² [[^]] <i>message-expression</i>
<i>conditional-component-part-list</i>	→	<i>simple-component-part-list</i> if <i>boolean-expr</i> {~ <i>simple-component-part-list</i> if <i>boolean-expr</i> }*
<i>simple-component-part-list</i>	→	<i>component-part</i> {\+ <i>component-part</i> }*
<i>simple-expr</i>	→	<i>message-expression</i> <i>primary</i>

A *message-expression* is a Smalltalk-style message expression and a *boolean-expr* is a Smalltalk-style message expression that returns true or false. A Smalltalk-style message expression consists of a receiver, a selector and zero or more arguments. If there are no arguments, the message is called a unary message. A binary message has a single argument following a selector consisting of one or two non-alphanumeric characters, the second of which may not be a minus sign. The third type of message is a keyword message. The selector consists of one or more keywords, each with its associated argument. A keyword consists of an identifier followed by a colon.

A *primary* is a Smalltalk-style primary which may be a variable name, a literal or a block. When a "<" symbol is immediately followed by the "[]" or "||" symbol, it denotes a quantification and the "<" symbol is not interpreted as a Smalltalk operator. If a ">" symbol appears as the first non-white-space character on a line, it denotes the end of a quantification construct.

No messages related to the Smalltalk-80 support for multiple processes may be used, since there is no concept of a process in a SLOOP program.

Cascaded message expressions are also not allowed. The motivation for this restriction was given in Chapter 4, Section 4.3.6.2.

A.8 Comments in a SLOOP program

Comments may be inserted anywhere in the program and are enclosed by double quotes.

² The braces around the [[^]]*variable* := *simple-expr* construct serve to identify the latter as a syntactic unit. This is needed because the '|' symbol has a higher precedence than the ':' symbol.

APPENDIX B

A SLOOP PROGRAM FOR A CALL CENTRE

B.1 Scope of the first level of refinement of the design

The first level of refinement is only concerned with normal behaviour; no error conditions are specified. This implies that a service user is always served once it is connected to the system; the possibility of aborting the connection (hanging up before the service has been completed) or rejecting the service request (e.g. due to unavailability of the relevant service providers) is not specified at this stage.

This is a cyclic system, which should therefore not terminate. However, there may be conditions under which the system may need to be shut down. System shutdown is not shown at this level of refinement.

The service providers may only be in the 'BUSY' or 'IDLE' states. At this level of refinement there is no `state` instance variable. The 'BUSY' and 'IDLE' status of a service provider is determined by checking whether it can accept the next service request or not. The 'RESTRICTED-IDLE', 'RESTRICTED-BUSY' and 'UNAVAILABLE' states are introduced during subsequent refinements. The additional states enable a service provider to go out of service gracefully and it ensures that once a service request is allocated to a service queue that it will be serviced.

(When a service provider is operating in a restricted mode, it means that it will continue to remain active until all the service requests already present in the service queue at the time when it changes to the restricted mode have been serviced. This indicates to the `scAllocator` that no further service requests should be accepted for a specific service queue if all the service providers servicing that particular category are operating in a restricted fashion.)

At this level of refinement the service user is not informed of the progress of the service request.

An object diagram of the Call Centre system is shown in Figure B-1. Figure B-2 illustrates the contents of the various packages in the system.

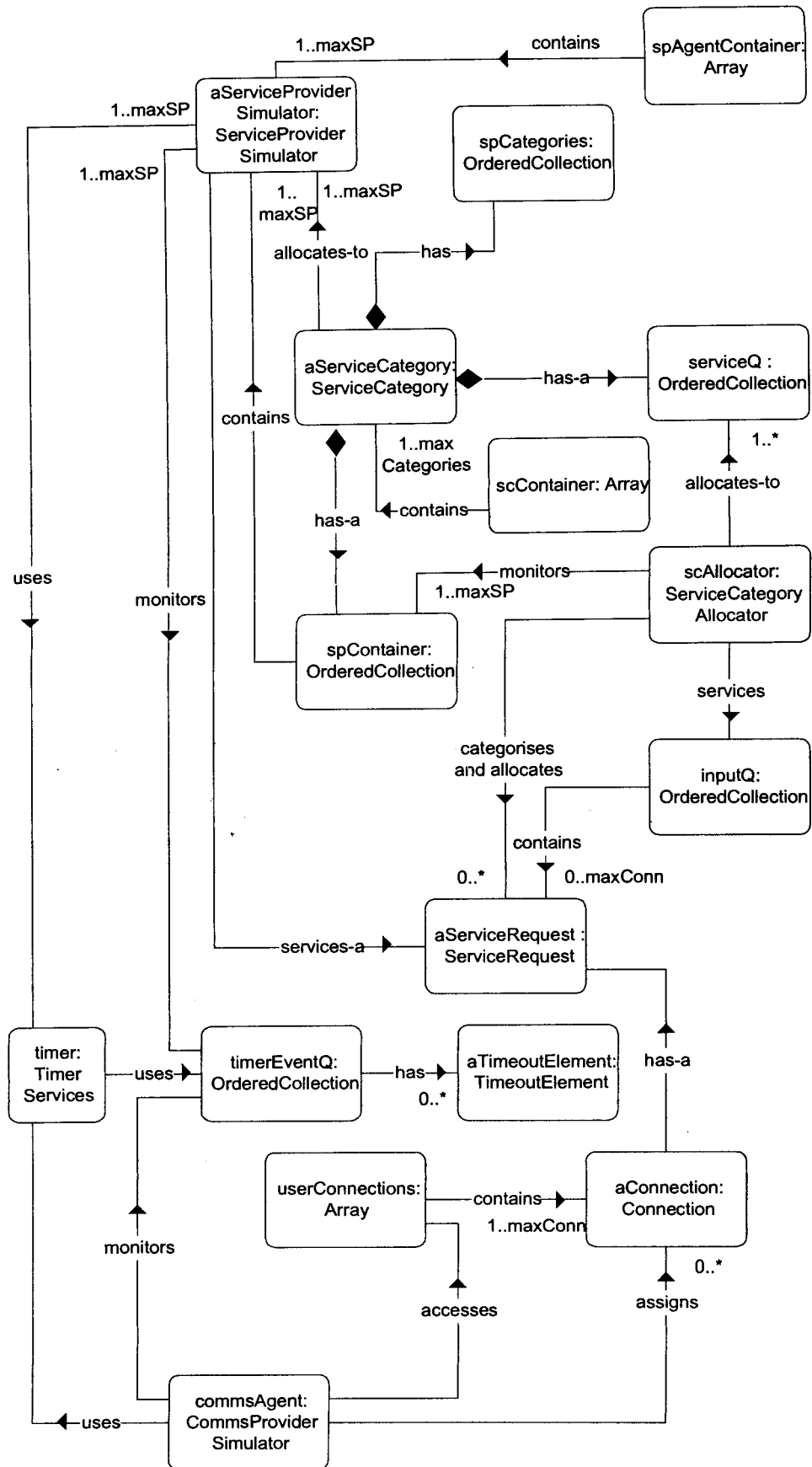


Figure B-1. Call centre object diagram.

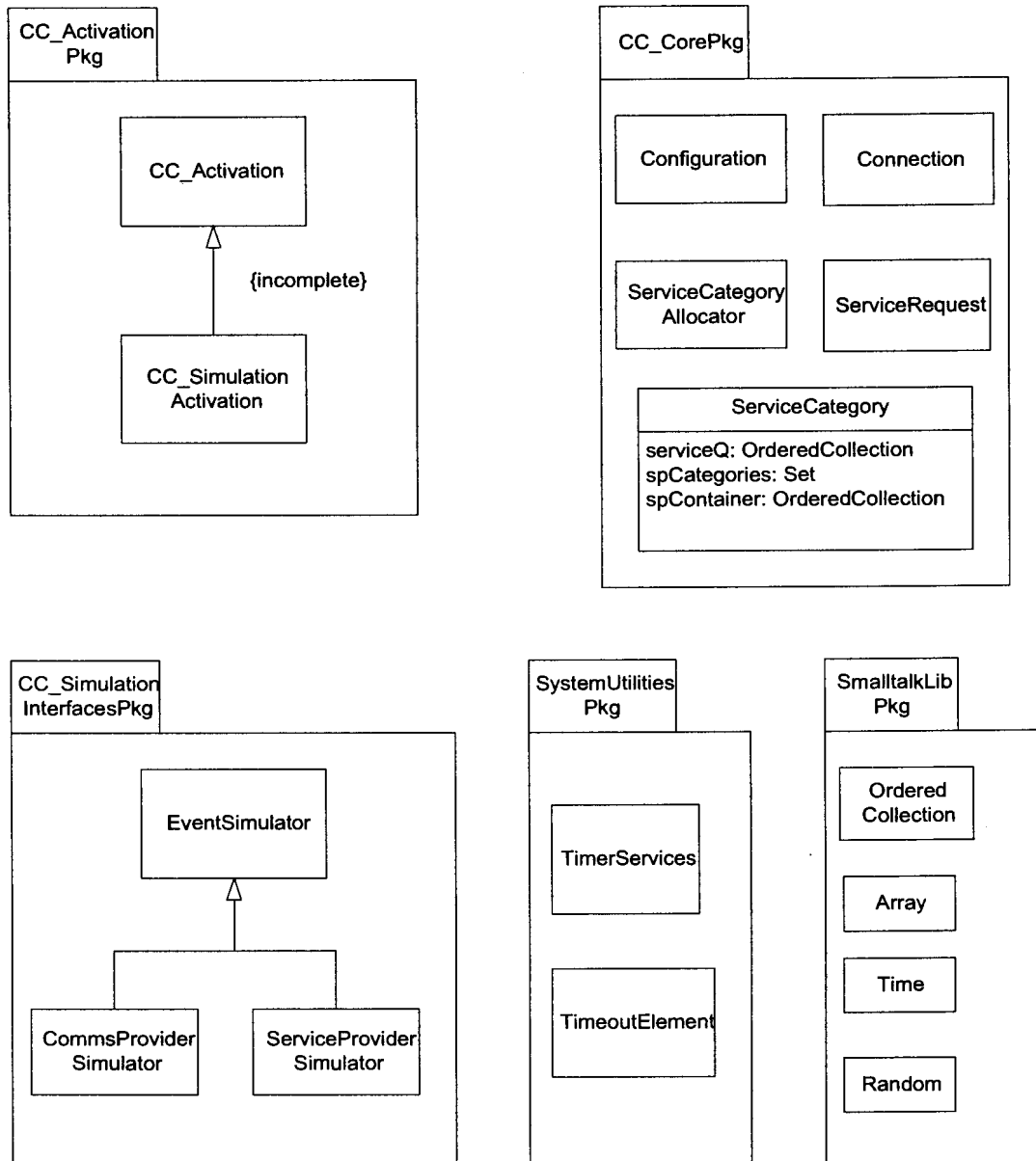


Figure B-2. Contents of the various call centre packages.

Full class descriptions are presented here for all non-Smalltalk library classes.

The property numbers are defined uniquely with respect to the class in which they appear. These numbers may appear in any order as long as they are unique. These numbers are independent of the numbers of the superclass properties. If a property overrides a property specified in an ancestor, the number of the ancestor property is used in the subclass followed by the name of the ancestor in brackets. When referring to a property that overrides a property in its superclass, the property identifier as described above must be followed by the words in *class-name*, where *class-name* is the name of the subclass.

In order to demonstrate that the concept of a package has no significance other than organizational, the classes in this Appendix are presented as individual classes and in any order.

B.2 The CC_Activation class

The `CC_Activation` class is an abstract superclass which leaves the activation of the **interface** classes up to the subclasses. It activates all the other classes in the system. The `CC_SimulationActivation` class overrides the methods related to the interface classes. The analysis level properties are not repeated here. They can be found in Chapter 5, Section 5.4.

```
class CC_Activation
"The CC Activation instance instantiates all the classes of the system
that need to present before the parallel methods start executing. It
also activates all the parallel methods required by the system."
```

```
superclass SmalltalkLibPkg::Object from SmalltalkLibRepository
```

```
instance variable names
```

```
"The following variables represent the objects that are not
instantiated as elements of a collection, i.e. they are not indexed
instance variables."
```

```
config
    "All configurable values (e.g. the maximum number of connections
    supported) are obtained via the config object."
commsAgent
    "The commsAgent handles the interface with the communication
    provider."
userConnections
    "This variable represents the collection of connections
    supported by the system."
inputQ
    "The inputQ models the FIFO way in which service requests are
    accepted by the system."
scAllocator
    "The scAllocator categorises the service requests and allocates
    them to the appropriate service queues."
scContainer
    "This variable represents the collection of service categories
    supported by the system."
spAgentContainer
    "The spAgentContainer contains all the service providers
    supported by the system."
timer
    "The timer object provides timer services to the other objects
    in the system."
timerEventQ
    "The timerEventQ is used to inform the requestors of the various
    timeouts that a requested timeout has occurred."
```

```
class macros
```

```
maxConn ≡ config maximumConnections
    "Number of simultaneous user connections supported"
[] maxCategories ≡ config maximumServiceCategories
    "Number of service categories supported"
[] maxSP ≡ config maximumServiceProviders
    "Number of service providers supported"
```

```
class properties
```

```
"These are the properties as identified during the analysis phase, as
well as the following design level clean behaviour invariants:"
```

```
invariant
```

```
config notNil ^
commsAgent notNil ^
userConnections notNil ^
```

```

< ∀ i where 1 ≤ i ≤ maxConn :: (userConnections at: i) notNil
> ^
inputQ notNil ^
scAllocator notNil ^
scContainer notNil ^
< ∀ j where 1 ≤ j ≤ maxCategories ::
  (scContainer at: j) notNil
> ^
spAgentContainer notNil ^
< ∀ k where 1 ≤ k ≤ maxSP :: (spAgentContainer at: k) notNil
> ^
timer notNil ^
timerEventQ notNil ^
maxConn > 0 ^
maxCategories > 0 ^
maxSP > 0 "DS2-01"
"Clean behaviour"
"The CC Activation class is an abstract class and should not be
instantiated"
<∀ anObject :: invariant anObject class ~~ CC_Activation
> "DS2-02"

```

instance methods

category private

message pattern initialize

method properties

"Total correctness"

true **results-in**

```

  methodReturnValue = self ^
  config notNil ^
  self postconditions: (#initManagement) ^
  commsAgent notNil ^
  self postconditions: (#initCommsAgent) ^
  userConnections notNil ^
  < ∀ i where 1 ≤ i ≤ maxConn :: (userConnections at: i )
  notNil
  > ^
  inputQ notNil ^
  scAllocator notNil ^
  self postconditions: (#initServiceCategoryAllocator) ^
  scContainer notNil ^
  < ∀ j where 1 ≤ j ≤ maxCategories :: (scContainer at: j)
  notNil
  > ^
  spAgentContainer notNil ^
  < ∀ k where 1 ≤ k ≤ maxSP :: (spAgentContainer at: k)
  notNil
  > ^
  timer notNil ^
  (timer class) postconditions:(#setup:)
  withArguments: #(config) ^
  timerEventQ notNil "DL1-02"

```

"Note that the receiver of the postconditions:withArguments: message is the expression (timer class) instead of SystemUtilitiesPkg::TimerServices. This is done to facilitate subclassing without violating the correctness properties. If the actual class name had been used here, then the property would no longer have been valid if a subclass of TimerServices

had been instantiated at this point. Recall that correctness properties must be preserved during subclassing."

"At this stage (i.e. before the subclass has completed the execution of its instance creation method) the class invariants do not need to hold yet, so it should be stated explicitly that once the predicate *self postconditions: (#initManagement)* holds, it continues to hold. That is a requirement, since many of the subsequent statements in the method depend on it. Similarly for the other **stable** properties listed below."

```
stable config notNil ^ self postconditions: (#initManagement)
                                                    "DS4-01"
"Note that self postconditions: (#initManagement) implies
that:
maxConn > 0 ^
maxCategories > 0 ^
maxSP > 0 "

stable userConnections notNil                "DS4-02"
stable scContainer                          "DS4-03"
stable spAgentContainer notNil             "DS4-04"
```

```
sequential
config := self initManagement
[] commsAgent := self initCommsAgent
[] userConnections := SmalltalkLibPkg:::Array new: maxConn
[] < [] i where 1 ≤ i ≤ maxConn :: userConnections at: i
    put: (self initConnection: i)
    >
[] inputQ := SmalltalkLibPkg:::OrderedCollection new: maxConn
[] scAllocator := self initServiceCategoryAllocator
[] scContainer := SmalltalkLibPkg:::Array
    new: maxCategories
[] < [] j where 1 ≤ j ≤ maxCategories :: scContainer at: j
    put: (CC_CorePkg:::ServiceCategory setup: config)
    >
[] spAgentContainer :=
    SmalltalkLibPkg:::Array new: maxSP
[] < [] k where 1 ≤ k ≤ maxSP :: spAgentContainer at: k
    put: (self initSPAgent)
    >
[] timer := SystemUtilitiesPkg:::TimerServices setup: config
[] timerEventQ := SmalltalkLibPkg:::OrderedCollection new
end-sequential
```

```
message pattern initManagement
method properties
"Total correctness"
true results-in
    methodReturnValue notNil ^
    (methodReturnValue class) postconditions:(#setup)
                                                    "DL1-03"
    "Again the explicit reference to a class name (in this
    case CC_CorePkg:::Configuration) is avoided in order to
    ensure that subclasses do not violate the correctness
    property."
sequential
^CC_CorePkg:::Configuration setup
end-sequential
```

```

message pattern initConnection: index
method properties
"Total correctness"
true results-in
    methodReturnValue notNil ^
    (methodReturnValue class) postconditions:(#setup:)
    withArguments: #(index) "DL1-04"
"Again the explicit reference to a class name (in this case
CC_CorePkg:::Connection) is avoided."
sequential
^CC_CorePkg:::Connection setup: index
end-sequential

message pattern initCommsAgent
method properties
"Total correctness"
true results-in methodReturnValue notNil "DL1-05"
sequential
self subclassResponsibility
end-sequential

message pattern initSPAgent
method properties
"Total correctness"
true results-in methodReturnValue notNil "DL1-06"
sequential
self subclassResponsibility
end-sequential

message pattern initServiceCategoryAllocator
method properties
"Total correctness"
true results-in methodReturnValue notNil "DL1-07"
sequential
^CC_CorePkg:::ServiceCategoryAllocator setup
end-sequential

category cyclic
message pattern p_activate
method properties
"These are the properties as identified during the analysis
phase."
"The p_activate method is only invoked once the CC_Activation
subclass has been instantiated. The class invariants of the
CC_Activation subclass that has been instantiated are therefore
guaranteed to hold before the p_activate method is executed.
Each statement executed by the p_activate method has to preserve
these invariants."

parallel
self p_executeCPAgent
"The parallel methods of the commsAgent are not invoked
directly, but rather via the p_executeCPAgent method of the
CC_Activation class."

[] timer p_runTimer: timerEventQ
"Activate the parallel methods of the timer object. The timer
parallel statements have the following functionality: Whenever a
timeout occurs, the TimeoutElement instance representing the
timeout is added to the end of the timerEventQ, which indicates
to the requestor that the specified timer has expired."

```



```
[] self p_categoriseAndAllocate
"The parallel methods of the sAllocator object are invoked via
the p_categoriseAndAllocate method of the CC_Activation class.
The sAllocator parallel statements have the following
functionality: Once a service request has been categorised, it
is removed from the inputQ and appended to the appropriate
serviceQ."

[] < [] j where 1≤j≤maxCategories :: (scContainer at: j)
    p_execute
>
"Activate the parallel methods of the ServiceCategory instances.
Their parallel statements have the following functionality: For
each service category the associated service queue and set of
service provider agents are monitored. If the service queue is
not empty and a service provider agent in the spSubset
associated with the service category is available to process a
new service request, the first element of the service queue is
removed and assigned to a service provider agent."

[] < [] i where 1≤i≤maxConn :: self p_executeConnection:
    (userConnections at: i)
>
"The p_executeConnection method of the CC_Activation class is
executed for each Connection instance in order to invoke the
parallel methods of the latter. The parallel statements of the
Connection instances have the following functionality: When a
connection has entered the 'TERMINATING' state, the
communication provider agent is requested to terminate the
connection. Once all the procedures have been completed to
terminate the connection, the connection and its associated
service request are reset to their initial states."

[] < [] k where 1≤k≤maxSP :: self p_executeSPAgent:
    (spAgentContainer at: k)
>
"The parallel methods of the service provider agents are not
invoked directly, but rather by executing the p_executeSPAgent
method of the CC_Activation class for each of the service
provider agents."

end-parallel

message pattern p_executeCPAgent
method properties
"These are the properties pertaining to the communication
provider interface as identified during the analysis phase."
parallel
self subclassResponsibility
end-parallel

message pattern p_executeSPAgent: spAgent
method properties
"These are the properties pertaining to the service provider
interface as identified during the analysis phase."
parallel
self subclassResponsibility
end-parallel

message pattern p_categoriseAndAllocate
method properties
"These are the properties pertaining to the service category
allocator as identified during the analysis phase."
```



parallel

scAllocator p_categorise: inputQ using: scContainer

"The scAllocator monitors the inputQ. If it is not empty, it enables the categorisation of the first element (a service request)."

[] scAllocator p_allocate: scContainer from: inputQ

"Once the service request has been categorised, the scAllocator removes it from the inputQ and appends it to the appropriate serviceQ."

end-parallel

message pattern p_executeConnection: aConnection

method properties

"These are the properties pertaining to the Connection class as identified during the analysis phase."

parallel

aConnection p_informCommsProvider: commsAgent

"When a connection has entered the 'TERMINATING' state, the communication provider agent is requested to terminate the connection."

[] aConnection p_doWrapUp

"Once all the procedures have been completed to terminate a connection, the connection and its associated service request are reset to their initial states."

end-parallel

B.3 The CC_SimulationActivation class

```

class CC_SimulationActivation
superclass CC_Activation
class properties
class methods
category instance creation
  message pattern setup
  method properties
  "Total correctness:
  After the statements in the initialize method have been executed
  the clean behaviour invariant of the CC_Activation class will
  hold, i.e. all the objects that should be created upon start-up
  will exist. This implies that the invariants of the classes
  instantiated by the CC_Activation class will also hold, as well
  as the correctness properties of their respective creation
  methods. All classes have to preserve their respective class
  invariants after initialisation."
  <∀ k where k ≥ 0 ::
  self instanceCount = k results-in
    self instanceCount = k + 1 ∧
    methodReturnValue notNil
  >
  sequential
  ^super new initialize
end-sequential
"DL1-01"

instance methods
category private
  message pattern initCommsAgent
  method properties
  "Total correctness"
  true results-in
    methodReturnValue notNil ∧
    CC_SimulationInterfacesPkg::CommsProviderSimulator
    postconditions: (#startSimulation)
    "DL1-05 (CC_Activation)"
  sequential
  ^CC_SimulationInterfacesPkg::CommsProviderSimulator
  startSimulation
end-sequential

  message pattern initSPAgent
  method properties
  "Total correctness"
  "The references to scContainer in the precondition of this
  method are required, because during initialization the
  ServiceProviderSimulator instance registers itself with each
  ServiceCategory instance that is serviced by a service provider
  category matching that of the current ServiceProviderSimulator
  instance."
  config notNil ∧ self postconditions: (#initManagement) ∧
  scContainer notNil ∧
  <∀ j where 1 ≤ j ≤ maxCategories :: (scContainer at: j) notNil
  > results-in
    methodReturnValue notNil ∧
    CC_SimulationInterfacesPkg::ServiceProviderSimulator
    postconditions: (#startSimulation:using:) withArguments:
    #(scContainer config)
    "DL1-6 (CC_Activation)"

```


sequential

```
^CC_SimulationInterfacesPkg::ServiceProviderSimulator
  startSimulation: scContainer using: config
```

end-sequential

category cyclic

message pattern p_executeCPAgent

method properties

"These are the properties pertaining to the communication provider interface as identified during the analysis phase."

parallel

```
commsAgent p_simulate: timer timeoutEventsIn: timerEventQ
```

```
[] commsAgent p_generateEvent: userConnections target: inputQ
```

"The commsAgent simulates the establishment of new connections at random intervals (within a configured range). A simulation timer is started after initialization and restarted each time after the establishment of a connection has been simulated. The latter is done by placing the service request associated with the new connection into the input queue. The commsAgent ensures that the capacity of maxConn connections per call centre is not exceeded, therefore a message is displayed indicating that all connections are busy if the maximum number of connections are currently assigned."

end-parallel

message pattern p_executeSPAgent: spAgent

method properties

"These are the properties pertaining to the service provider interface as identified during the analysis phase."

parallel

```
spAgent p_simulate: timer timeoutEventsIn: timerEventQ
```

"When a service request has been assigned to a service provider simulator, the latter simulates the time it takes to service the service request by starting a random timer. When this timer expires, it represents the completion of the service."

```
[] spAgent p_generateEvent
```

"When the service provider has completed the service, it indicates that the connection should be terminated."

```
[] spAgent p_updateCategoryIndex: scContainer
```

"Update the index into the categoriesServed collection if the serviceQ of the current category being served by this spAgent is empty."

end-parallel

B.4 The Configuration class

class Configuration

"The purpose of this class is to ensure that the following parameters are configured:

Uses defaults to set the maximum number of connections, service categories and service providers. It also facilitates the configuration of the maximum allowable timeout value, the service request category names supported by the system, the service provider category names supported by the system and the mapping of service request to service provider category names. Subclasses may allow the operator to specify other values."

superclass Object

instance variable names

maximumConnections

"The maximum number of connections supported by the system."

maximumServiceCategories

"The number of service categories supported by the system."

maximumServiceProviders

"The number of service providers supported by the system."

maximumAllowableTimeout

"The maximum allowable timeout that may be requested by any object in the system."

srCategoryNames

"The collection of service request category names supported by the system."

spCategoryNames

"The collection of service provider category names supported by the system."

srToSpCategoryMap

"The mapping of service request categories to service provider categories."

categoriesAssigned

"This variable is used to keep track of the number of service request category names that have already been assigned. When a ServiceCategory instance is created and initialized, it obtains the name of the service category that it supports from the Configuration instance (via the assignSRCategory method). Each service request category name may only be assigned once (in order to ensure that each ServiceCategory instance will support a unique service request category)."

class properties

$\langle \forall (t, u, v, w) \text{ where}$

$t > 0 \wedge u > 0 \wedge v > 0 \wedge w > 0 ::$

invariant

maximumConnections = t \wedge

maximumServiceCategories = u \wedge

maximumServiceProviders = v \wedge

maximumAllowableTimeout = w

\rangle

"DS2-01"

"The values of each of the maximumConnections, maximumServiceCategories, maximumServiceProviders and maximumAllowableTimeout instance variables are invariant and always greater than zero."

```

<∀ u where u > 0 ::
invariant
    srCategoryNames notNil ∧ srCategoryNames size = u
>
    "DS3-01"
"The number of service request category names that are configured is equal to
maximumServiceCategories."

```

```

<∀ ( anSRCategoryNameX, anSRCategoryNameY) where
srCategoryNames includes: anSRCategoryNameX ∧
srCategoryNames includes: anSRCategoryNameY ::
invariant
    anSRCategoryNameX ~~ anSRCategoryNameY
>
    "DS3-02"
"Each configured service request category name is unique."

```

```

invariant
    spCategoryNames notNil ∧ ¬spCategoryNames isEmpty
    "DS3-03"
"At least one service provider category name is configured."

```

```

<∀ ( anSPCategoryNameX, anSPCategoryNameY) where
spCategoryNames includes: anSPCategoryNameX ∧
spCategoryNames includes: anSPCategoryNameY ::
invariant
    anSPCategoryNameX ~~ anSPCategoryNameY
>
    "DS3-04"

```

class methods

```

category instance creation
message pattern setup
method properties
    "A Configuration instance is created and initialized. The new
instance is returned"

    "Total correctness"
<∀ k where k ≥ 0 ::
self instanceCount = k results-in
    self instanceCount = k + 1 ∧
    methodReturnValue notNil
>
    "DL1-01"
sequential
^super new configure
end-sequential

```

instance methods

```

category private
message pattern configure
method properties
    "Upon completion of the configure method, the
maximumConnections, maximumServiceCategories, maximumService=
Providers and maximumAllowableTimeout instance variables will
each have a value greater than zero, the srCategoryNames and
spCategoryNames collections will have been created, the number
of elements in the srCategoryNames collection will be equal to
maximumServiceCategories, there will be at least one element in
the spCategoryNames collection, the srToSpCategoryMap will have
been created, the number of mappings in this collection will be
equal to maximumServiceCategories and the categoriesAssigned
variable will have the value zero."

```



```

"Total correctness"
<∀ ( t, u, v, w) where
t > 0 ∧ u > 0 ∧ v > 0 ∧ w > 0 ::
true results-in
    maximumConnections = t ∧
    maximumServiceCategories = u ∧
    maximumServiceProviders = v ∧
    maximumAllowableTimeout = w ∧
    srCategoryNames notNil ∧ spCategoryNames notNil ∧
    srCategoryNames size = u ∧
    ¬ spCategoryNames isEmpty ∧
    srToSpCategoryMap notNil ∧
    srToSpCategoryMap size = u ∧
    categoriesAssigned = 0
>
sequential
[] maximumConnections := 8
    "Maximum number of simultaneous user connections"
[] maximumServiceCategories := 1
    "Maximum number of service categories"
[] maximumServiceProviders := 3
    "Maximum number of service providers"
[] maximumAllowableTimeout := 5
    "Maximum allowable timeout"
[] srCategoryNames := SmalltalkLibPkg::OrderedCollection
    new: maximumServiceCategories
[] srCategoryNames addLast: 'Default Service Request category'
[] spCategoryNames := SmalltalkLibPkg::OrderedCollection
    new: maximumServiceProviders
"Multiple service providers may belong to the same service
provider category, but it is also possible that each service
provider could belong to a different service provider category.
The maximum size is therefore used when spCategoryNames is
created."
[] spCategoryNames addLast: 'Default Service Provider category'
[] srToSpCategoryMap = Dictionary new
[] srToSpCategoryMap at: 'Default Service Request category'
    put: spCategoryNames
[] categoriesAssigned := 0
end-sequential

```

"DL1-02"

category accessing

```

message pattern maximumConnections
method properties
"Total correctness"
true results-in methodReturnValue = maximumConnections "DL1-03"
sequential
^maximumConnections
end-sequential

```

```

message pattern maximumServiceCategories
method properties
"Total correctness"
true results-in methodReturnValue = maximumServiceCategories
sequential
^maximumServiceCategories
end-sequential

```

"DL1-04"

```

message pattern maximumServiceProviders
method properties
"Total correctness"
true results-in methodReturnValue = maximumServiceProviders
"DL1-05"

```

```

sequential
^maximumServiceProviders
end-sequential

```

```

message pattern maximumAllowableTimeout
method properties
"Total correctness"
true results-in methodReturnValue = maximumAllowableTimeout
"DL1-06"

```

```

sequential
^maximumAllowableTimeout
end-sequential

```

```

message pattern getSPCategories: srCategory
method properties
"The srCategory passed as a parameter is used as the index into
the srToSpCategoryMap object in order to extract the collection
of Service Provider Categories associated with the srCategory."
"Total correctness"
true results-in
    methodReturnValue = srToSpCategoryMap at: srCategory
"DL1-07"

```

```

sequential
^srToSpCategoryMap at: srCategory
end-sequential

```

category modifying

```

message pattern assignSRCategory
method properties
"Returns a unique service request category (each service request
category is only assigned once)"
"Total correctness"
<∀ x where 0 ≤ x < maximumServiceCategories ::
categoriesAssigned = x results-in
    categoriesAssigned = x + 1 ∧
    methodReturnValue = srCategoryNames at: (x + 1)
>
"DL1-08"

```

```

sequential
categoriesAssigned := categoriesAssigned + 1
[] ^srCategoryNames at: categoriesAssigned
end-sequential

```

```

message pattern assignSPCategory
method properties
"Returns a service provider category (each service provider
category may be assigned multiple times). Subclasses may use
various algorithms to assign service provider categories"

```



"Total correctness"

```
<∀ x where spCategoryNames includes: x ::  
true results-in  
    methodReturnValue = x  
>
```

"DL1-09"

```
sequential  
^spCategoryNames first  
end-sequential
```

B.5 The EventSimulator class

The EventSimulator class is an abstract class. It is responsible for starting a timer if one is required. It also detects the expiry of the timer. The subclasses of EventSimulator are responsible for determining **when** a timer is required and also for **generating the events** resulting from the expiry of the timers.

```

class EventSimulator
superclass Object from SmalltalkLibRepository
instance variable names
rand
    "This variable refers to an instance of the Random class from
    the Smalltalk library. The instance is created when the
    EventSimulator subclass is instantiated. The instance of the
    Random class maintains a seed from which the next random number
    is generated. The random number is used to start a timer with a
    random value."
newEventRequired
    "When the value is equal to true it means that a new event is
    required. Once the variable has been set to true, a random
    timer will be started at some point afterwards. When the timer
    is started, newEventRequired is set to false. It is the
    responsibility of the subclass to set this variable to true when
    a new event is required, since each subclass will have its own
    conditions for requiring a new event. Once the timer expires,
    an event will be generated, as will be described in the comments
    section of the generatingEvent variable."
currentRandomTimeoutValue
    "This variable contains the value of the random timeout
    currently being requested. The purpose of this variable is to
    provide a mechanism for referencing the current timeout value in
    the correctness arguments. Note that the SLOOP statements could
    therefore have been rewritten without this variable while still
    providing the same functionality. However, in that case it
    would not have been possible to formalise certain correctness
    properties (such as DL1-04)."
generatingEvent
    "The value is equal to true if the timer has expired and an
    event has to be generated, otherwise it is equal to false. The
    subclass sets this variable to false at the time when the event
    is generated. The actual event that is generated is also the
    responsibility of the subclass, since each subclass will generate
    a different type of event."
timerOutstanding
    "This variable is set to true when a timer is started and it is
    set to false when a timeoutElement is removed from the
    timerEventQ (i.e. when an expired timer has been processed).
    The purpose of this variable is to provide a mechanism for
    reasoning about the uniqueness of outstanding timers in the
    EventSimulator class. In this class only one timer requested by
    the EventSimulator may be outstanding at a time. The
    timerOutstanding variable is used in the preconditions of the
    startRandomTimer:withMaximum: method as well as in the
    postconditions of the resetTimerExpired: method. If subclasses
    need to support multiple simultaneous timers, then the
    preconditions of the startRandomTimer:withMaximum: method need
    to be weakened and the postconditions of the resetTimerExpired:
    method need to be strengthened. Since the purpose of the
    timerOutstanding variable is to facilitate correctness
    reasoning, the SLOOP statements could have been rewritten
    without this variable while still providing the same
    functionality."

```

```

timerId
    "This variable contains the identifier of the timer currently
    being requested."

class properties

    "Liveness"
    "When a simulation event is required, a simulation timer is eventually started."
    "AL2-01"

    "Liveness"
    "If a simulator timer expires, the simulator eventually has to generate an event."
    "AL2-02"

    "Clean behaviour"
    <∀ anObject ::
        invariant anObject class ~~ EventSimulator
    >
    "DS2-01"
    "The EventSimulator class is an abstract class and should not be instantiated"

    "Clean behaviour"
    invariant rand notNil ^ rand class = Random "DS2-02"
    "Once rand has been initialized to refer to an instance of the Random class, it is never
    set to nil while the instance of the EventSimulator subclass exists."
    "It is therefore possible for the EventSimulator subclass instance to send messages
    to rand at any stage after initialization."

    "Clean behaviour"
    "The currentRandomTimeout value is always within the range specified by the
    precondition of the start:id:for: method of the TimerServices class."
    "DS2-03"

    "Global invariant"
    "All outstanding timers requested by an EventSimulator subclass instance are
    identified uniquely with respect to the requestor."
    "DS3-01"
    "Thus, all the timers requested by this requestor that are
    currently running or that are in the timerEventQ are uniquely
    identified with respect to the requestor."

instance methods
category private
    message pattern initialize
    "Creates an instance of class Random and sets newEventRequired,
    generatingEvent and timerOutstanding to false. It also sets
    currentRandomTimeoutValue to 1 so that the class invariant
    referring to it will hold after instance creation and
    initialization have been completed."
    method properties
    "Total correctness"
    true results-in methodReturnValue = self ^
    rand notNil ^ newEventRequired = false ^
    currentRandomTimeoutValue = 1 ^
    generatingEvent = false ^
    timerOutstanding = false
    "DL1-01"

    sequential
    rand := SmalltalkLibPkg::Random new
    [] newEventRequired := false
    [] currentRandomTimeoutValue := 1

```



```
[] generatingEvent := false
>[] timerOutstanding := false
end-sequential
```

category accessing

```
message pattern nextRandomNumber: maximumValue
method properties
"Returns the next random number between 1 and maximumValue
inclusive"
"Total correctness"
true results-in
    methodReturnValue ≥ 1 ∧ methodReturnValue ≤ maximumValue
sequential
^ (rand next * maximumValue) truncated + 1
end-sequential
"DL1-02"
```

category testing

```
message pattern timerExpired: timerEventQ
method properties
"Returns true if the timerEventQ contains an element of which
the requestor == self, otherwise it returns false. This method
needs to be overridden when multiple simultaneous timers may be
originated by the same requestor. In that case the identifier
which uniquely identifies the timer with respect to the
requestor has to match as well."
"Total correctness"
true results-in methodReturnValue =
    (timerEventQ detect: [:each |
        each timeoutRequestor == self ]
        ifNone: [nil]) notNil
sequential
^ (timerEventQ detect: [:each | each timeoutRequestor == self]
    ifNone: [nil]) notNil
end-sequential
"DL1-03"
```

category modifying

```
message pattern startRandomTimer: aTimerServices
    withMaximum: maximumValue
"Start a timer with a random value within the range between 1
and maximumValue. When the resulting start:id:for: message is
sent to the TimerServices instance, a reference to the requestor
(in this case the EventSimulator subclass instance) as well as
an identifier are passed as parameters. The combination of the
reference to the requestor and the identifier ensures that each
timer request can be identified uniquely within the system.
This facilitates the correlation of the subsequent timeout
notifications with the timer requests.
```

In the EventSimulator class only one timer is outstanding at a time for a specific requestor, i.e. `¬timerOutstanding` is a precondition for starting a new timer for a specific instance of an EventSimulator subclass. Since the timers initiated by a specific EventSimulator subclass instance do not run concurrently, these timers can all have an identifier of 1.

If a subclass requires multiple concurrent timers, unique values must be allocated to the corresponding identifiers. The `startRandomTimer:withMaximum:` method therefore needs to be overridden in order to achieve this. The total correctness property of the modified method also needs to be updated, viz. a **disjunction** needs to be added to the precondition to state that the proposed identifier of any new timer requested by that

EventSimulator subclass instance should not match any identifier of any other outstanding timer requested by that EventSimulator subclass instance. Thus, the precondition has to be **weakened**. In that case the value of timerOutstanding will no longer be relevant."

method properties

"Total correctness"

```
¬timerOutstanding results-in methodReturnValue = self ∧
  self postconditions: (#nextRandomNumber:)
  withArguments: #(maximumValue) ∧
  aTimerServices postconditions: (#start:id:for:)
  withArguments: #(self timerId currentRandomTimeoutValue) ∧
  timerOutstanding "DL1-04"
```

sequential

```
currentRandomTimeoutValue :=
  (self nextRandomNumber: maximumValue)
```

```
[] timerId := 1
```

```
[] aTimerServices start: self id: timerId for:
  currentRandomTimeoutValue
```

```
[] timerOutstanding := true
```

end-sequential

message pattern resetTimerExpired: timerEventQ

method properties

"Removes the first timeoutElement in timerEventQ where the requestor matches the receiver."

"Total correctness"

```
<∃ expiredTimeout where expiredTimeout ==
  (timerEventQ detect: [:each | each timeoutRequestor == self]
  ifNone: [nil]) ::
  expiredTimeout notNil results-in
  methodReturnValue = self ∧
  (timerEventQ includes: expiredTimeout) not ∧
  timerOutstanding not
```

>

"DL1-05"

sequential

```
timerEventQ remove:
```

```
(timerEventQ detect: [:each | each timeoutRequestor == self])
```

```
[] timerOutstanding := false
```

end-sequential

category cyclic

message pattern p_simulate: aTimerServices timeoutEventsIn:
timerEventQ

"If a new event is required, start a random timer, the expiry of which will cause an event to be initiated."

method properties

"This method ensures that properties **DS2-03**, **AL2-01** and **AL2-02** are satisfied by the EventSimulator class."

"Clean behaviour"

```
invariant currentRandomTimeoutValue > 0 ∧
  currentRandomTimeoutValue ≤ aTimerServices
  maximumTimeout "DS2-03"
```

"A timeout requested by the EventSimulator subclass instance is always within the range that ensures that the precondition of the start:id:for: method of the

TimerServices class is met when the EventSimulator subclass instance invokes that method."

```

"Precedence"
newEventRequired ensures
    self postconditions: (#startRandomTimer:withMaximum:)
    withArguments:
        #(aTimerServices (aTimerServices maximumTimeout))
        ^ -newEventRequired "DP1-01"
"When newEventRequired is true, it ensures that a simulation timer is started and
newEventRequired becomes false."

"Precedence"
self timerExpired: timerEventQ ensures
    generatingEvent ^
    self postconditions: (#resetTimerExpired:)
    withArguments: #(timerEventQ) "DP1-02"
"When a simulation timer expires, it ensures that generatingEvent becomes true."

parallel
self startRandomTimer: aTimerServices withMaximum:
(aTimerServices maximumTimeout) \+
newEventRequired := false
    if newEventRequired
[] generatingEvent := true \+
self resetTimerExpired: timerEventQ
    if self timerExpired: timerEventQ
end-parallel

```

B.6 The CommsProviderSimulator class

The CommsProviderSimulator class simulates the actions of the communication provider. This class is replaced by the CommunicationProviderAgent class when the call centre interacts with the actual communication provider instead of simulating its actions.

The CommsProviderSimulator class is a subclass of EventSimulator. The newEventRequired instance variable is set to true as part of the initialization procedures. That results in the starting of a new timer. Once the timer has expired, an event is generated. A new connection is established if one is available and the associated service request is added to the end of the input queue. The timeout is ignored if all the connections are busy. This would correspond to a busy signal being received by the service user in an actual implementation.

```

class CommsProviderSimulator
superclass EventSimulator from ApplicationsRepository
class properties
"The communication provider agent constructs the serviceRequest
associated with the connection based on information received from the
communication provider. For example, in the one case the service user
identification information and the type of service required will be
received. In another case no information will be received (perhaps
because it is not relevant, e.g. when a directory enquiry is made).
The simulation puts no information into the service request."

class methods
category instance creation
  message pattern startSimulation
  method properties

  "The initialize method of the superclass sets newEventRequired
  to false. This method sends the initialize message to its
  superclass and immediately sets newEventRequired to true. The
  values of the other instance variables mentioned in the total
  correctness property of the initialize method of the superclass
  remain unchanged."
  "Total correctness"
  <∀ k where k ≥ 0 ::
  self instanceCount = k results-in
    self instanceCount = k + 1 ∧
    methodReturnValue notNil
  >
  sequential
  ^ (super new initialize) moreInit
  end-sequential
  "DL1-01"

instance methods
category private
  message pattern moreInit
  "Sets newEventRequired to true (which was set to false in the
  initialization routine of the superclass)."  

  method properties
  "Total correctness"
  true results-in methodReturnValue = self ∧
    newEventRequired = true
  "DL1-02"
  "This property ensures that property AS3-01, which was
  identified during the analysis phase, is achieved. Property
  AS3-01 specifies that instance creation results in a
  communication provider simulator event being required."

```

```
sequential
newEventRequired := true
end-sequential
```

```
category accessing
message pattern getIdleConnection: userConnections
"Returns the first connection that is idle"
method properties
"Total correctness"
true results-in methodReturnValue =
    userConnections detect: [:each | each isIdle]
    ifNone: [nil] "DL1-03"
sequential
^ userConnections detect: [:each | each isIdle] ifNone: [nil]
end-sequential
```

```
category modifying
message pattern terminate: aConnection cause: reason
method properties
"Inform the communication provider that the connection has
terminated."
"Total correctness"
true results-in methodReturnValue = self "DL1-04"
sequential
Transcript show: 'Connection'
[] Transcript show: (aConnection connectionIndex printString)
[] Transcript show: 'has terminated with cause'
[] Transcript show: reason
end-sequential
```

```
category cyclic
message pattern p_generateEvent: userConnections
    target: inputQ
method macros
idleConnection ≡ self getIdleConnection: userConnections
method properties
"Simulate an event from the communication provider."

"Safe liveness"
generatingEvent ^ idleConnection notNil ^
¬(inputQ includes:(idleConnection serviceRequest)) ensures
    ¬generatingEvent ^ newEventRequired ^
    inputQ last = (idleConnection serviceRequest) ^
    idleConnection postconditions: (#assign) "AP1-01"
"If an event has to be generated and the maximum number of connections have not yet
been established, the communication provider simulator ensures that a new connection
is established, the associated service request is appended to the input queue and a new
communication provider simulator event is again required."

"Safe liveness"
generatingEvent ^ idleConnection isNil ensures
    ¬generatingEvent ^ newEventRequired "AP1-02"
"If an event has to be generated and the maximum number of connections have already
been established, the communication provider simulator ensures that the event is
cancelled and a new communication provider simulator event is again required."

parallel
inputQ addLast: (idleConnection serviceRequest) \+
idleConnection assign
    if generatingEvent and: [idleConnection notNil] ~
```

```
Transcript show: 'All connections busy'  
    if generatingEvent and: [idleConnection isNil]  
    || newEventRequired := true \+  
generatingEvent := false  
    if generatingEvent  
end-parallel
```

B.7 The Connection class

```

class Connection
superclass Object from SmalltalkLibRepository
instance variable names
state
    "The state of the connection"
terminatingReason
    "The reason why the connection is being terminated."
serviceRequest
    "The service request associated with the connection."
currentHandlerInformed
    "This flag is used when a connection has to be terminated. It
    indicates whether the communication provider has been informed
    of the termination of the connection."
connectionIndex
    "The index of this connection into the userConnections array."

```

class properties

"Note that there are no safety properties specifying the allowed values of the state instance variable. There are also no safety properties specifying the allowed state transitions. The reason for this is to avoid overspecification, i.e. it avoids restricting subclasses to certain specified values. Recall that preconditions may not be strengthened and postconditions may not be weakened during subclassing."

class methods

```

category instance creation
message pattern setup: indexOfConnection
method properties
    "Total correctness"
    <∀ k where k ≥ 0 ::
    self instanceCount = k results-in
        self instanceCount = k + 1 ∧
        methodReturnValue notNil
    >
sequential
super new initialize: indexOfConnection
end-sequential

```

"DL1-01"

instance methods

```

category private
message pattern initialize: indexOfConnection
method properties
    "Total correctness"
    true results-in methodReturnValue = self ∧ state = 'IDLE' ∧
    serviceRequest notNil ∧ currentHandlerInformed = false ∧
    connectionIndex = indexOfConnection
sequential
state := 'IDLE'
[] serviceRequest := CC_CorePkg::ServiceRequest setup: self
[] currentHandlerInformed := false
[] connectionIndex := indexOfConnection
end-sequential

```

"DL1-11"

```

category accessing
  message pattern terminatingReason
  method properties
  "Total correctness"
  true results-in methodReturnValue = terminatingReason "DL1-02"
  sequential
  ^ terminatingReason
  end-sequential

  message pattern connectionIndex
  method properties
  "Total correctness"
  true results-in methodReturnValue = connectionIndex "DL1-03"
  sequential
  ^ connectionIndex
  end-sequential

  message pattern serviceRequest
  method properties
  "Total correctness"
  true results-in methodReturnValue = serviceRequest "DL1-04"
  sequential
  ^ serviceRequest
  end-sequential

category testing
  message pattern isIdle
  method properties
  "Total correctness"
  true results-in methodReturnValue = (state = 'IDLE') "DL1-05"
  sequential
  ^ state = 'IDLE'
  end-sequential

  message pattern isTerminating
  method properties
  "Total correctness"
  true results-in methodReturnValue = (state = 'TERMINATING')
  "DL1-06"
  sequential
  ^ state = 'TERMINATING'
  end-sequential

category modifying
  message pattern assign
  method properties
  "Total correctness"
  state = 'IDLE' results-in methodReturnValue = self ^ state =
  'CONNECTED' "DL1-07"
  sequential
  state := 'CONNECTED'
  if state = 'IDLE'
  end-sequential

```



```

message pattern terminate: reason
method properties
"Total correctness"
state = 'CONNECTED' results-in
    methodReturnValue = self ^
    state = 'TERMINATING' ^ terminatingReason = reason
                                                                    "DL1-08"
"Total correctness"
state = 'TERMINATING' results-in methodReturnValue = self
                                                                    "DL1-09"

"This allows for terminate collision."

"Total correctness"
state = 'IDLE' results-in methodReturnValue = self
                                                                    "DL1-10"
"This ensures that the transition from 'IDLE' to 'TERMINATING'
is not possible"
sequential
terminatingReason := reason \+
state := 'TERMINATING'
    if state = 'CONNECTED'
end-sequential

category cyclic
message pattern p_informCommsProvider: commsAgent
method properties
"Safe liveness"
state = 'TERMINATING' ^ terminatingReason = 'completed' ^
~currentHandlerInformed ensures
    commsAgent postconditions: (#terminate:cause:)
    withArguments: #(self terminatingReason) ^
    currentHandlerInformed
                                                                    "DP1-01"
parallel
commsAgent terminate: self cause: terminatingReason \+
currentHandlerInformed := true
    if state = 'TERMINATING' and:
        [(terminatingReason = 'completed')
        and: [currentHandlerInformed not]]
end-parallel

message pattern p_doWrapUp
method properties
"Safe liveness"
currentHandlerInformed ensures
state = 'IDLE' ^ serviceRequest postconditions: (#reset) ^
~currentHandlerInformed
                                                                    "DP1-02"
parallel
state := 'IDLE' \+
serviceRequest reset \+
currentHandlerInformed := false
    if currentHandlerInformed
end-parallel

```

B.8 The ServiceCategoryAllocator class

```

class ServiceCategoryAllocator
superclass Object from SmalltalkLibRepository
instance variable categorising
    "The categorising variable is used as a flag to indicate whether
    the categorisation of the service request at the head of the
    inputQ has been initiated or not."
class properties
"Safe liveness"
<∀ aServiceRequest where inputQ includes: aServiceRequest ::
    inputQ first = aServiceRequest ∧ ¬categorising ensures
    inputQ first = aServiceRequest ∧ categorising
>
"Safe liveness"
"DP1-01"
<∀ aServiceRequest where inputQ includes: aServiceRequest ::
    inputQ first = aServiceRequest ∧ categorising ensures
    < ∃ aServiceQueue where
        (scContainer detects: [:each | each serviceQ = aServiceQueue]
        ifNone: [nil]) notNil ::
        aServiceQueue includes: aServiceRequest)
    > ∧ ¬(inputQ includes: aServiceRequest) ∧ ¬(categorising)
>
"DP1-02"

class methods
category instance creation
message pattern setup
method properties
    "Total correctness"
<∀ k where k ≥ 0 ::
    self instanceCount = k results-in
        self instanceCount = k + 1 ∧
        methodReturnValue notNil
>
"DL1-01"
sequential
super new initialize
end-sequential

instance methods
category private
message pattern initialize
method properties
    "Total correctness"
true results-in methodReturnValue = self ∧ ¬categorising
"DL1-02"
sequential
categorising := false
end-sequential

category modifying
message pattern categoriseServiceRequest: serviceRequest
    using: scContainer
    "When the service request does not contain any categorisation
    data, it is categorised as belonging to the first service
    category in scContainer. This is the default behaviour which
    facilitates usage of this class without further subclassing if
    only one service category is supported by the system. This
    method needs to be reimplemented in the subclasses if multiple
    service categories are supported. In that case the default
    category is only used if the categorisation data is not

```

provided, otherwise the service request is categorised according to the data provided by the service user (e.g. via the IVR)."

method properties

"Total correctness"

true results-in

```

methodReturnValue = self ^
serviceRequest serviceRequestCategory notNil ^
(scContainer detects:
[:each | each serviceQCategory =
serviceRequest serviceRequestCategory] ifNone: [nil]))
notNil
"DL1-03"

```

sequential

```

serviceRequest serviceRequestCategory:
(scContainer first) serviceQCategory

```

end-sequential

message pattern assignToSQ: serviceRequest using: scContainer

method macros

```

match ≡ scContainer detect: [:each | each serviceQCategory =
serviceRequest serviceRequestCategory] ifNone: [nil]

```

method properties

"Further refinements would override this method to include error conditions. For example, the service request would be rejected if the service providers in the service provider container were all operating in the restricted mode."

"Total correctness"

```

serviceRequest serviceQ isNil ^
serviceRequest serviceRequestCategory notNil ^ match notNil

```

results-in

```

methodReturnValue = self ^
serviceRequest serviceQ notNil ^
serviceRequest serviceRequestCategory notNil ^
(match serviceQ) includes: serviceRequest
"DL1-04"

```

sequential

```

serviceRequest serviceQ: (match serviceQ) \+
(match serviceQ) addLast: serviceRequest
if match notNil

```

end-sequential

category cyclic

message pattern p_categorise: inputQ using: scContainer

method properties

"Safe liveness"

```

¬(inputQ isEmpty) ^ ¬categorising until
categorising ^
self postconditions: (#categoriseServiceRequest:using:)
withArguments: #((inputQ first) scContainer)
"DP1-03"

```

parallel

```

categorising := true \+
self categoriseServiceRequest: (inputQ first)
using: scContainer
if inputQ isEmpty not and: [categorising not]

```

end-parallel

message pattern p_allocate: scContainer from: inputQ

method macros

```

serviceRequest ≡ inputQ first
if inputQ isEmpty not ~
nil
if inputQ isEmpty

```



method properties

"Safe liveness"

<∀ aServiceRequest **where**

¬(inputQ isEmpty) ∧ inputQ first == aServiceRequest ::

 aServiceRequest serviceRequestCategory notNil ∧

 aServiceRequest serviceQ isNil ∧

 < ∃ aServiceCategory **where**

 scContainer includes: aServiceCategory ::

 aServiceCategory serviceQCategory =

 aServiceRequest serviceRequestCategory

 >

ensures

 self postconditions: (#assignToSQ:using:)

 withArguments: #(aServiceRequest scContainer) ∧

 ¬categorising ∧

 ¬(inputQ includes: aServiceRequest)

>

"DP1-04"

parallel

 self assignToSQ: serviceRequest using: scContainer \+

 categorising := false \+

 inputQ removeFirst

 if serviceRequest notNil and:

 [serviceRequest serviceRequestCategory notNil and:

 [serviceRequest serviceQ isNil]]

end-parallel

B.9 The ServiceRequest class

```

class ServiceRequest
  superclass Object from SmalltalkLibRepository
  instance variable names
  serviceQ
    "The service queue to which the service request has been
    assigned."
  serviceRequestCategory
    "The category of this service request."
  connection
    "The connection associated with this service request."
  serviceProvider
    "The service provider to which this service request has been
    assigned."
  categorisationData
    "The data which is used to categorise this service request."

  class properties

  class methods
  category Instance creation
  message pattern setup: aConnection
  method properties
  "Total correctness"
  <∀ k where k ≥ 0 ::
  self instanceCount = k results-in
    self instanceCount = k + 1 ∧
    methodReturnValue notNil
  >
  sequential
  ^super new initialize: aConnection
  end-sequential
  "DL1-01"

  instance methods
  category private
  message pattern initialize: aConnection
  method properties
  true results-in methodReturnValue = self ∧
    connection = aConnection ∧
    self postconditions: (#reset)
  sequential
  connection := aConnection
  [] self reset
  end-sequential
  "DL1-02"

  category accessing
  message pattern serviceQ
  method properties
  "Total correctness"
  true results-in methodReturnValue = serviceQ
  sequential
  ^ serviceQ
  end-sequential
  "DL1-03"

```

```

message pattern serviceRequestCategory
method properties
"Total correctness"
true results-in methodReturnValue = serviceRequestCategory
"DL1-04"

```

```

sequential
^ serviceRequestCategory
end-sequential

```

```

message pattern connection
method properties
"Total correctness"
true results-in methodReturnValue = connection
"DL1-05"

```

```

sequential
^ connection
end-sequential

```

```

message pattern serviceProvider
method properties
"Total correctness"
true results-in methodReturnValue = serviceProvider
"DL1-06"

```

```

sequential
^ serviceProvider
end-sequential

```

```

message pattern categorisationData
method properties
"Total correctness"
true results-in methodReturnValue = categorisationData
"DL1-07"

```

```

sequential
^ categorisationData
end-sequential

```

category modifying

```

message pattern serviceQ: sq
method properties
"Total correctness"
true results-in methodReturnValue = self ^ serviceQ = sq
"DL1-08"

```

```

sequential
serviceQ := sq
end-sequential

```

```

message pattern serviceRequestCategory: srCategory
method properties
"Total correctness"
true results-in methodReturnValue = self ^
    serviceRequestCategory = srCategory
"DL1-09"

```

```

sequential
serviceRequestCategory := srCategory
end-sequential

```

```

message pattern serviceProvider: sp
method properties
"Total correctness"
true results-in methodReturnValue = self ^
    serviceProvider = sp
"DL1-10"

```

```

sequential
serviceProvider := sp
end-sequential

```

```
message pattern categorisationData: newData
method properties
"Total correctness"
true results-in methodReturnValue = self ^
    categorisationData = newData                                "DL1-11"
sequential
categorisationData := newData
end-sequential

message pattern reset
method properties
"Total correctness"
true results-in methodReturnValue = self ^ serviceQ isNil ^
    serviceRequestCategory isNil ^ serviceProvider isNil ^
    categorisationData isNil                                    "DL1-12"
sequential
serviceQ := nil
[] serviceRequestCategory := nil
[] serviceProvider := nil
[] categorisationData := nil
end-sequential
```

B.10 The ServiceCategory class

Subclasses may override the *canAssignSR* and *assignToSP* methods to implement algorithms to assign service requests to service providers as required by individual applications. The algorithm implemented for the *ServiceQueueCategory* class is to assign the service request to the first available service provider in the container.

In anticipation of a refinement which might require that service providers be allocated in a round robin fashion, *OrderedCollection* rather than *Set* is used for the service provider container component of the *ServiceCategory* class. The main difference between *OrderedCollection* and *Set* is the fact that the elements of the former are ordered and those of the latter are not. By ordering the service providers within such a container, it makes it possible to implement an algorithm which will ensure that all service providers are utilised. For example, if service requests of a specific category arrive slowly enough that the second service request is only considered once the service provider has finished serving the first, then it could happen that the service requests are always allocated to the same service provider. An *OrderedCollection* implementation could assist in ensuring that the service requests are allocated to the service providers in a round robin fashion.

```

class ServiceCategory
superclass Object

instance variable names
serviceQCategory
    "The category of the service requests enqueued in the serviceQ"
serviceQ
    "The FIFO queue containing service requests matching
    serviceQCategory"
spCategories
    "The collection of service provider categories that apply to the
    serviceQ"
spSubset
    "The collection of service providers that may service the
    serviceQ"

class properties

class methods
category instance creation
message pattern setup: config
method properties
    "Total correctness"
    <∀ k where k ≥ 0 ::
    self instanceCount = k ∧ config notNil results-in
        self instanceCount = k + 1 ∧
        methodReturnValue notNil
    >
    sequential
    ^super new initialize: config
    end-sequential
instance methods
category private
message pattern initialize: config
method macros
    maxConn ≡ config maximumConnections
  
```

"DL1-01"


```

maxSP = config maximumServiceProviders

method properties
"Total correctness"
config notNil results-in
  methodReturnValue = self ^
  serviceQ notNil ^ serviceQCategory notNil ^
  spSubset notNil ^ spCategories notNil ^
  self postconditions: (#registerSPCategories:
  withArguments: #(config) "DL1-02"
sequential
serviceQ := SmalltalkLibPkg::OrderedCollection new: maxConn
[] serviceQCategory := config assignSRCategory
[] spSubset :=
  SmalltalkLibPkg::OrderedCollection new: maxSP
[] spCategories := SmalltalkLibPkg::OrderedCollection new
[] self registerSPCategories: config
end-sequential

message pattern registerSPCategories: config
method properties
"Total correctness"
config notNil ^ serviceQCategory notNil results-in
  methodReturnValue = self ^ ¬spCategories isEmpty "DL1-03"
sequential
spCategories addAll:
  (config getSPCategories: serviceQCategory)
end-sequential

category testing
message pattern servicedBy: spCategory
method properties
"Total correctness"
spCategory notNil results-in
  methodReturnValue = (spCategories includes: spCategory)
"DL1-04"
sequential
^spCategories includes: spCategory
end-sequential

message pattern canAssignSR
method properties
"Total correctness"
true results-in methodReturnValue =
  (spSubset detect:
  [:each | each canAcceptNextSR: serviceQCategory]
  ifNone: [nil] ) notNil "DL1-05"
sequential
^ ( (spSubset detect:
  [:each | each canAcceptNextSR: serviceQCategory]
  ifNone: [nil]) notNil)
end-sequential

category accessing
message pattern spSubset
method properties
"Total correctness"
true results-in methodReturnValue = spSubset "DL1-06"

```

```
sequential
^spSubset
end-sequential
```

```
message pattern serviceQCategory
method properties
"Total correctness"
true results-in methodReturnValue = serviceQCategory    "DL1-07"
sequential
^serviceQCategory
end-sequential
```

```
message pattern serviceQ
method properties
"Total correctness"
true results-in methodReturnValue = serviceQ            "DL1-08"
sequential
^serviceQ
end-sequential
```

```
category modifying
message pattern addSP: anSP
method properties
"Total correctness"
anSP notNil results-in
    methodReturnValue = self ^ spSubset includes: anSP    "DL1-09"
sequential
spSubset addLast: anSP
end-sequential
```

```
message pattern assignToSP: sr
method macros
availableServiceProvider ≡
    spSubset detect: [:each | each canAcceptNextSR:
        serviceQCategory]
method properties
"Total correctness"
sr notNil ^ availableServiceProvider notNil results-in
methodReturnValue = self ^
sr postconditions: (#serviceProvider:)
    withArguments: #(availableServiceProvider) ^
availableServiceProvider
    postconditions: (#processServiceRequest:)
    withArguments: #(sr)    "DL1-10"
sequential
sr serviceProvider: availableServiceProvider
[] availableServiceProvider processServiceRequest: sr
end-sequential
```

```
category cyclic
message pattern p_execute
method properties
"Safe liveness"
serviceQ isEmpty not ^ self canAssignSR ensures
self postconditions: (#assignToSP:) withArguments:
    #((serviceQ first)) ^
serviceQ postconditions: (#removeFirst)    "DP1-01"
```



```
parallel  
self assignToSP: (serviceQ first) \+  
serviceQ removeFirst  
    if serviceQ isEmpty not and: [self canAssignSR]  
end-parallel
```

B.11 The TimerServices class

The TimerServices class allows its clients to request that timers of specified durations be started. The timer resolution is in seconds and the maximum timeout value is denoted by `maximumTimeout`. The TimerServices class uses a **circular** array called `timeoutCollection` to implement the timers. Each position in the array represents one second. The object maintains an index into the array. This index is called `currentTick` and is advanced every second. Thus, the entry in the array which will be reached x seconds from the current moment can be calculated using the value of `currentTick`, the size of the array and the value of x .

Each entry in the array is an ordered collection of `TimeoutElement` instances. When the TimerServices class receives a request to start a timer, it creates a `TimeoutElement` instance to represent that timer and enters it into the relevant collection of `TimeoutElement` instances, i.e. it enters it into the collection that will be reached after x seconds, where x is the timeout value specified for the timer. The `TimeoutElement` instance contains the timeout identifier. It also contains other attributes that may be used to obtain information about the timeout.

When a timer expires, its associated `TimeoutElement` instance is removed from the relevant collection in the circular array and it is added to `timerEventQ`. The latter is inspected by the clients of TimerServices in order to determine whether their requested timers have timed out yet.

The `currentTick` variable is used to calculate the read and write indices and is updated every second. When the TimerServices instance is created, `currentTick` is initialized to zero. Every second it is incremented modulo the size of the array. The size of the array is a function of the maximum timeout value.

Thus, when a TimerServices instance receives a request to start a timer, it has to determine which list of `TimeoutElement` instances will be processed after $(\text{duration} + 1)$ clock ticks. It is necessary to add the extra one in order to ensure that a timer does not expire prematurely. For example, if a timer for one second is started at the time when half a second of the `currentTick` period has already passed, then the new timer will expire after only half a second. By adding the extra one, a timer of one second could take up to two seconds to expire, but it will at least run for one second.

The write index is therefore calculated as follows:

```
writeIndex :=
((currentTick + duration + 1) \% (timeoutCollection size)) + 1
```

The requested timeout value (represented by `duration`) plus one is added to the value of `currentTick`. The addition is performed modulo the size of the array, which is `maximumTimeout + 2`. The size of the array is explained as follows: Since one is always added to the `duration` in order to ensure that a timer does not expire prematurely, the maximum timeout value is actually `maximumTimeout + 1`. Another one has to be added to ensure that the read and write indices never have the same value. After the modulo addition has been performed, the `writeIndex` has to be adjusted by one, since SLOOP array indices start at one, not zero. The `readIndex` has the value `currentTick` plus one. It is also adjusted by one in order to compensate for the fact that the array indices start at one. Once the `writeIndex` has been calculated, the `TimeoutElement` instance is added to the end of the First In First Out list that is stored at that position in the array.

Figure B-3(a) illustrates where a new entry would be added for the following values of the various variables:

maximumTimeout = 5
 duration = 5
 currentTick = 3

Figure B-3(b) shows the result if the currentTick value had been 0 at the time of the timer request.

Note that the currentTick value ranges from 0 to 6 inclusive in the above example, whereas readIndex and writeIndex may range from 1 to 7 inclusive.

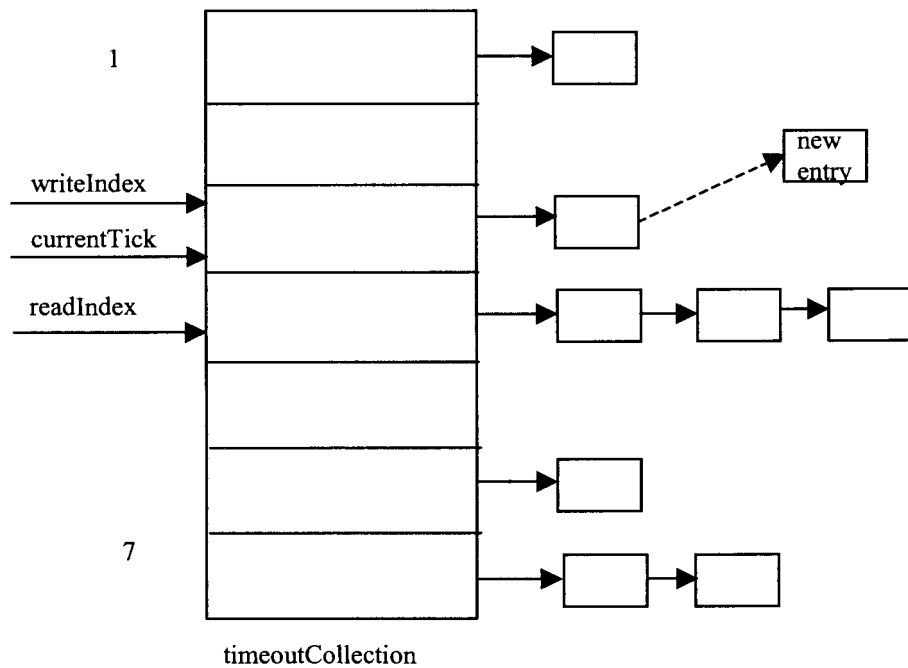


Figure B-3(a). Structures used by aTimerServices (currentTick = 3).

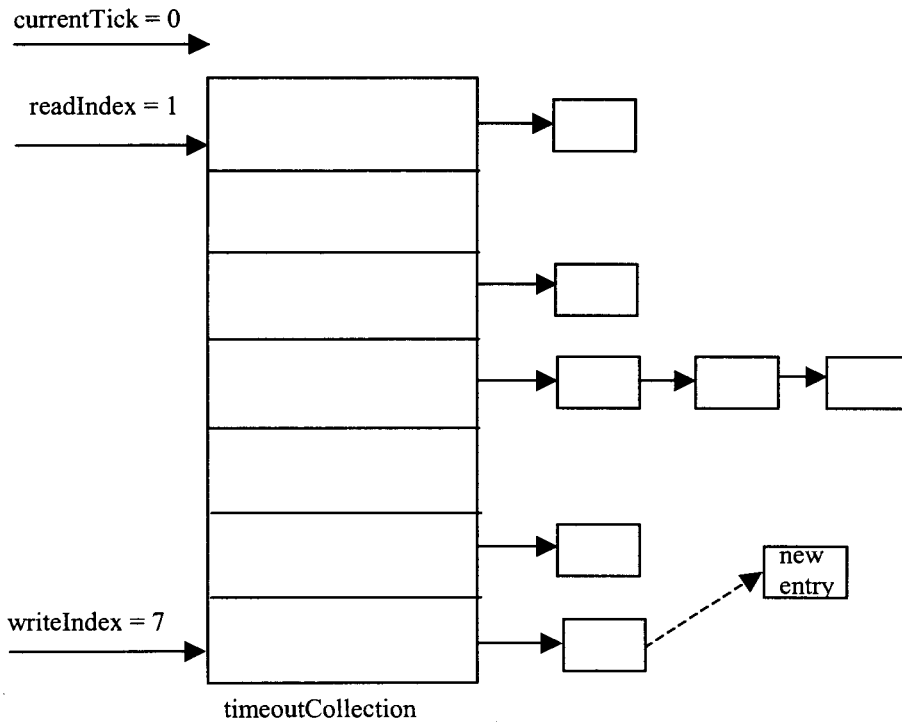


Figure B-3(b). Structures used by aTimerServices (currentTick = 0).

The value of the currentTick variable is updated whenever one second has passed since its last update and all the entries at the current readIndex have been processed. This is reflected by one of the parallel statements of the TimerServices class (it is executed infinitely often):

```

lastTime := currentTime \+
currentTick := (currentTick + 1) \\\ (timeoutCollection size)
  if difference ≥ 1 and: [currentTimeoutElement isNil]
  
```

This statement uses two *macro-variables*, viz. difference and currentTimeoutElement. They receive their values in the *macro-section* of the method containing the parallel statement, as shown below (the *macro-variable* readIndex is used in the definition of currentTimeoutElement and therefore has to be defined prior to its usage):

```

  readIndex := currentTick + 1
  [] difference ≡ currentTime - lastTime
    if (currentTime - lastTime) ≥ 0 ~
      currentTime + (86400 - lastTime)
    if (currentTime - lastTime) < 0
  
```

```
[] currentTimeoutElement :=  
  (timeoutCollection at: readIndex) first  
  if (timeoutCollection at: readIndex) isEmpty not ~  
    currentTimeoutElement := nil  
  if (timeoutCollection at: readIndex) isEmpty
```

One of the other parallel statements of the `TimerServices` class ensures that the `currentTime` variable is updated on a regular basis:

```
currentTime := SmalltalkLibPkg::Time now asSeconds
```

Evaluation of the `Time now asSeconds` expression yields the number of seconds since midnight. The calculation of difference takes the rollover at midnight into account in the *macros-section* shown above.

The `TimerServices` class is now presented in the SLOOP notation.

```

class TimerServices
superclass Object from SmalltalkLibRepository
instance variable names
maximumTimeout
    "The maximum timeout value that may be requested"
timeoutCollection
    "A circular array. Each element comprises a list of
    TimeoutElement instances"
currentTick
    "It points to a position in the timeoutCollection array. It is
    used to calculate the read and write indices."
currentTime
    "The most recent time (in number of seconds since midnight)
    obtained from the system."
lastTime
    "The time (in number of seconds since midnight) when the
    currentTick was last updated."

class properties
"Global invariant: when a timeout is indicated to the
requestor, then a period greater than or equal to the value
specified by the requestor has expired."
invariant
<∀ aTimeoutElement where
    timerEventQ includes: aTimeoutElement::
    aTimeoutElement timerExpired
>
"DS3-01"

<∀ aTimeoutElement where
    <∃ i where 1 ≤ i ≤ (timeoutCollection size) ::
    (timeoutCollection at: i) includes: aTimeoutElement
    > ::
    -aTimeoutElement timerServicesCompleted leads-to
        aTimeoutElement timerServicesCompleted
    >
    "DL2-01 (TimerServices)"
"Once a timer has been started, i.e. it is present in one of the lists associated with
    timeoutCollection, the TimerServices instance will eventually complete its
    responsibilities regarding the timer (i.e. the timer will either be stopped or the
    TimerServices instance will indicate its expiry to the requestor of the timer)."
    "This property refers to the timerServicesCompleted method
    rather than the timerExpired method in order to make provision
    for subclasses that may allow a timer to be aborted."

class methods
category Instance creation
message pattern setup: config
method properties
    "Total correctness"
    <∀ k where k ≥ 0 ::
    self instanceCount = k ∧ config notNil results-in
        self instanceCount = k + 1 ∧
        methodReturnValue notNil
    >
    "DL1-01"
sequential
    ^super new initialize: config
end-sequential

```


instance methods

category private

message pattern initialize: config

method properties

"Total correctness"

true **results-in** timeoutCollection notNil ^

maximumTimeout = config maximumAllowableTimeout ^

< \forall i **where** $1 \leq i \leq$ timeoutCollection size ::

(timeoutCollection at: i) notNil

> ^

currentTick = 0 ^ currentTime notNil ^ lastTime notNil

"DL1-02"

sequential

maximumTimeout := config maximumAllowableTimeout

[] timeoutCollection :=

SmalltalkLibPkg::Array new: (maximumTimeout + 2)

[] < [] i **where** $1 \leq i \leq$ timeoutCollection size ::

timeoutCollection at: i put: (OrderedCollection new)

>

[] currentTick := 0

[] currentTime := SmalltalkLibPkg::Time now asSeconds

[] lastTime := currentTime

end-sequential

category accessing

message pattern maximumTimeout

method properties

"Total correctness"

true **results-in** methodReturnValue = maximumTimeout

"DL1-03"

sequential

^maximumTimeout

end-sequential

message pattern isTimerRunningFor: requestor with: identifier

method properties

"Total correctness"

true **results-in** methodReturnValue = (found notNil) ^

found **detects**

< \forall i **where** ($1 \leq i \leq$ timeoutCollection size) ::

< \exists aTimeoutElement **where**

(timeoutCollection at: i) includes: aTimeoutElement ::

aTimeoutElement timeoutRequestor = requestor ^

aTimeoutElement timeoutIdentifier = identifier

>

>

"DL1-04"

sequential

found := nil

[] < [] i **where** ($1 \leq i \leq$ timeoutCollection size) ::

found := (timeoutCollection at: i) detect:

[:each | each timeoutRequestor = requestor and:

[each timeoutIdentifier = identifier]] ifNone: [nil]

if found isNil

>

[] ^ found notNil

end-sequential

category modifying

```

message pattern start: requestor id: identifier for: duration
method macros
  writeIndex ≡
    ((currentTick + duration + 1) \ (timeoutCollection size))
    + 1
    "This is because the array index starts at 1, not 0"

method properties
"Clean behaviour"
invariant 1 ≤ writeIndex ∧
  writeIndex ≤ timeoutCollection size "DS3-02"

invariant writeIndex ≠ readIndex "DS3-03"

"Total correctness"
0 < duration ∧ duration ≤ maximumTimeout results-in
  methodReturnValue = self ∧
  nextElement class = TimeoutElement ∧
  (timeoutCollection at: writeIndex) includes: nextElement ∧
  TimeoutElement postconditions: (#setup:id:for:)
  withArguments: #(requestor identifier duration) "DL1-05"

sequential
nextElement :=
SystemUtilitiesPkg::TimeoutElement setup: requestor
id: identifier for: duration
  if 0 < duration ∧ duration ≤ maximumTimeout
[] (timeoutCollection at: writeIndex) addLast: nextElement
  if 0 < duration ∧ duration ≤ maximumTimeout
end-sequential

```

category cyclic

```

message pattern p_runTimer: timerEventQ
method macros
  readIndex ≡ currentTick + 1
  "This is because the array index starts at 1, not 0"

[] difference ≡ currentTime - lastTime
  if (currentTime - lastTime) ≥ 0 ~
  currentTime + (86400 - lastTime)
  if (currentTime - lastTime) < 0

[] currentTimeoutElement ≡
  (timeoutCollection at: readIndex) first
  if (timeoutCollection at: readIndex) isEmpty not ~
  nil
  if (timeoutCollection at: readIndex) isEmpty

method properties
"Intermittent assertion"
< ∀ y where 0 ≤ y < (timeoutCollection size) ::
currentTick = y ∧ difference ≥ 1 ∧
(currentTimeoutElement isNil leads-to
  currentTick = (y + 1) \ (timeoutCollection size)
> "DL2-02"
"Thus, all the entries at the current timeout position have to
be processed before the entries at the next position are
processed. The granularity of the timer is ≥ 1 second"

```

```

"Safe liveness"
<∀aTimeoutElement where
aTimeoutElement = currentTimeoutElement::
aTimeoutElement = notNil ensures
    ¬((timeoutCollection at: readIndex) includes:
    aTimeoutElement) ^
    timerEventQ includes: aTimeoutElement ^
    aTimeoutElement timerServicesCompleted
>
"DP1-01"

parallel
    currentTime := SmalltalkLibPkg::Time now asSeconds
[] lastTime := currentTime \+
    currentTick := (currentTick + 1) \\ (timeoutCollection size)
    if difference ≥ 1 and: [currentTimeElement isNil]
[] timerEventQ addLast: currentTimeoutElement \+
    currentTimeoutElement updateEndTime \+
    currentTimeoutElement timerServicesCompleted: true \+
    (timeoutCollection at: readIndex) removeFirst
    if currentTimeoutElement notNil
end-parallel

```

Note that the macro-variable `writeIndex` cannot be defined in the class-macros section, because it refers to the pseudo-variable `duration`.

The clients of the `TimerServices` class are not restricted to removing only the first element of the `timerEventQ`. That ensures that each client will receive its timeout information even if other clients misbehave. Each client must remove the `TimeoutElement` instance from the `timerEventQ` as it processes it, otherwise it will process the same element multiple times.

The advantage of having a `timerEventQ` is that the interface between the `TimerServices` instance and its clients is very loosely coupled. Alternatively the `TimerServices` instance can invoke a client method when a timer expires, but in that case it is necessary to reserve the client together with the `TimerServices` instance when the timeout is processed.

The reason for defining `timerEventQ` as a peer class rather than as part of an aggregation (i.e. within `TimerServices`) is to allow for more parallelism. This way it is not necessary to reserve `timerEventQ` as well whenever a timer method is executed.

B.12 The TimeoutElement class

The TimeoutElement class is now presented using the SLOOP notation:

```

class TimeoutElement
superclass Object from SmalltalkLibRepository
instance variable names
startTime
    "The time when this timeout was started"
endTime
    "The time when this timeout expired."
timeoutRequestor
    "The requestor of this timeout."
timeoutIdentifier
    "The identifier of this timeout. It is unique with respect to
    the requestor."
requestedDuration
    "The requested duration of this timeout."
timerServicesCompleted
    "This flag indicates whether the TimerServices instance has
    completed its tasks regarding this timeout. This flag is used
    rather than checking the endTime and startTime in order to be
    able to cater for the case where the timeout is aborted."
class macros
    currentTime  $\equiv$  SmalltalkLibPkg::Time totalSeconds
    "currentTime contains the total number of seconds since January
    1, 1901."

class properties
    invariant    endTime  $\geq$  startTime                "DS3-01"
    self getCurrentDuration = 0 unless
        self getCurrentDuration > 0                    "DS4-01"

class methods
category instance creation
    message pattern setup: requestor id: identifier for: duration
    method properties
    "Total correctness"
    < $\forall$  k where k  $\geq$  0 ::
    self instanceCount = k results-in
        self instanceCount = k + 1  $\wedge$ 
        methodReturnValue notNil
    >
    "DL1-01"
    sequential
    ^super new initialize: requestor id: identifier for: duration
    end-sequential

instance methods
category private
    message pattern initialize: requestor id: identifier for:
    duration
    method properties
    "Total correctness"
    true results-in methodReturnValue = self  $\wedge$ 
        timeoutRequestor = requestor  $\wedge$ 
        timeoutIdentifier = identifier  $\wedge$ 
        startTime notNil  $\wedge$  endTime notNil  $\wedge$ 
        requestedDuration = duration  $\wedge$   $\neg$ timerServicesCompleted
    "DL1-02"

```

```

sequential
timeoutRequestor := requestor
[] timeoutIdentifier := identifier
[] startTime := currentTime
[] endTime := currentTime
[] requestedDuration := duration
[] timerServicesCompleted := false
end-sequential

```

category accessing

```

message pattern timeoutRequestor
method properties
"Total correctness"
true results-in methodReturnValue = timeoutRequestor "DL1-03"
sequential
^timeoutRequestor
end-sequential

```

```

message pattern timeoutIdentifier
method properties
"Total correctness"
true results-in methodReturnValue = timeoutIdentifier "DL1-04"
sequential
^timeoutIdentifier
end-sequential

```

```

message pattern startTime
method properties
"Total correctness"
true results-in methodReturnValue = startTime "DL1-05"
sequential
^startTime
end-sequential

```

```

message pattern requestedDuration
method properties
"Total correctness"
true results-in methodReturnValue = requestedDuration "DL1-06"
sequential
^requestedDuration
end-sequential

```

```

message pattern getCurrentDuration
method properties
"Total correctness"
true results-in methodReturnValue = (currentTime - startTime) "DL1-07"
sequential
^(currentTime - startTime)
end-sequential

```

```

message pattern getTimeoutDuration
method properties
"Total correctness"
true results-in methodReturnValue = (endTime - startTime) "DL1-08"
sequential
^(endTime - startTime)
end-sequential

```

```

category testing
  message pattern timerExpired
  method properties
    " Total correctness "
    true results-in methodReturnValue =
      (self getCurrentDuration - requestedDuration ≥ 0)
                                                    "DL1-09"

  sequential
    ^(self getCurrentDuration - requestedDuration ≥ 0)
  end-sequential

  message pattern timerServicesCompleted
  method properties
    "Total correctness"
    true results-in methodReturnValue = timerServicesCompleted
                                                    "DL1-10"

  sequential
    ^timerServicesCompleted
  end-sequential

category modifying
  message pattern updateEndTime
  method properties
    "Total correctness"
    true results-in methodReturnValue = self ^
      endTime = currentTime
                                                    "DL1-11"

  sequential
    endTime := currentTime
  end-sequential

  message pattern timerServicesCompleted: newValue
  method properties
    "Total correctness"
    true results-in methodReturnValue = self ^
      timerServicesCompleted = newValue
                                                    "DL1-12"

  sequential
    timerServicesCompleted := newValue
  end-sequential

```

B.13 The ServiceProviderSimulator class

class ServiceProviderSimulator

superclass EventSimulator **from** ApplicationsRepository

instance variable names

serviceRequest

"This variable refers to the service request currently being serviced by the service provider simulator. Note that the reference to the ServiceRequest instance is passed to the simulator as a parameter, i.e. the ServiceRequest instance is not created by the ServiceProviderSimulator instance and therefore does not form part of it."

serviceProviderCategory

"This variable contains the name of the service provider category to which the service provider simulator belongs."

categoriesServed

"This is an ordered collection containing the names of the service request categories serviced by this service provider. The purpose of this array is to facilitate a round robin servicing scheme of these categories. That prevents starvation of a specific service category."

nrOfCategoriesServed

"This variable contains the number of service request categories serviced by this service provider. It is used in the calculation when the categoryIndex is updated."

categoryIndex

"This variable is used as index into the categoriesServed collection. It is used to determine the next service request category to be serviced by this service provider. It is incremented modulo nrOfCategoriesServed. Its values range from 0 to nrOfCategoriesServed - 1"

class properties

$\langle \forall \text{ categoryIndex where } \text{categoryIndex} \geq 0 \wedge$

$\text{categoryIndex} \leq \text{nrCategoriesServed} - 1 ::$

invariant serviceRequest notNil $\Rightarrow \neg \text{self canAcceptNextSR:}$
(categoriesServed at: (categoryIndex + 1))

\rangle

"AS3-01"

"A service provider simulator services a single service request at a time."

"If a service request is currently assigned to the simulator, no other service request from any of the categories being served by this simulator will be served by the latter."

serviceRequest isNil $\wedge \neg \text{newEventRequired unless}$
serviceRequest notNil $\wedge \text{newEventRequired}$

"AS4-01"

"When a new service request is assigned to the service provider simulator then a new service provider simulator event is required."

Note: The parent class, viz. EventSimulator, contains a parallel method which monitors the value of newEventRequired. If it detects that newEventRequired is true, it starts a timer and sets newEventRequired to false."

serviceRequest notNil $\wedge \neg \text{newEventRequired unless}$
serviceRequest isNil $\wedge \neg \text{newEventRequired}$

"AS4-02"

"If a service request has been assigned to the service provider simulator and newEventRequired is false, then newEventRequired remains false while the service request is still assigned to the service provider simulator."

"This has the effect that this simulator will not start another timer before the servicing of the current service request has been completed."

```
generatingEvent ^ serviceRequest notNil ensures
  (serviceRequest connection)
  postconditions: (#terminate:)
  withArguments: #('completed') ^ serviceRequest isNil ^
  -generatingEvent "AP1-01"
```

"If a service provider simulator has to generate an event, it ensures that the service provider simulator terminates the connection currently associated with it and becomes available to service a new service request."

"Note: The parent class, viz. EventSimulator, contains a parallel method which sets generatingEvent to true when a timer has expired."

```
<∀ aServiceRequest where serviceRequest = aServiceRequest ::
  serviceRequest = aServiceRequest ensures
  (serviceRequest connection)
  postconditions: (#terminate:)
  withArguments: #('completed') ^ serviceRequest isNil
> "AP1-02"
```

"A service request remains assigned to a service provider simulator until the latter completes the service and terminates the connection."

```
invariant categoryIndex ≥ 0 ^
  categoryIndex < nrOfCategoriesServed "DS2-01"
```

"The categoryIndex is always greater than or equal to zero and less than nrOfCategoriesServed."

```
invariant categoriesServed notNil ^
  categoriesServed class = OrderedCollection "DS2-02"
```

"Once categoriesServed has been initialized to refer to an instance of the OrderedCollection class, it is never set to nil while the ServiceProviderSimulator instance exists."

```
<∀ categoryIndex where 0 ≤ categoryIndex ^
  categoryIndex < nrOfCategoriesServed ::
  ¬(self canAcceptNextSR:
  (categoriesServed at: (categoryIndex + 1)))
leads-to
  self canAcceptNextSR:
  (categoriesServed at: (categoryIndex + 1))
> "DL2-01"
```

"For any service category serviced by the service provider simulator, the service provider simulator will eventually be able to service a request from that service category."

class methods

category instance creation
message pattern startSimulation: scContainer using:
 aConfiguration
method properties
"The initialize method of the superclass sets newEventRequired to false. The startSimulation:using: method invokes the initialize message of its superclass and initialises the ServiceProviderSimulator-specific instance variables. The properties of the initialize method of the superclass hold."
"Total correctness"
true **results-in** methodReturnValue notNil ^
 super postconditions: (#initialize) ^
 serviceRequest isNil ^
 serviceProviderCategory notNil ^
 categoriesServed notNil ^
 nrOfCategoriesServed ≥ 0 ^
 categoryIndex ≥ 0
"AL1-01"
"Instance creation results in the initialization of the instance variables of the ServiceProviderSimulator class and its superclasses."
sequential
^(super new initialize) moreInit: scContainer using:
aConfiguration
end-sequential

instance methods

category private
message pattern moreInit: scContainer using: aConfiguration
"Initializes the ServiceProviderSimulator instance"
method properties
"Total correctness"
true **results-in** methodReturnValue = self ^
 serviceRequest isNil ^
 aConfiguration postconditions: (#assignSPCategory) ^
 serviceProviderCategory notNil ^
 categoriesServed notNil ^
 self postconditions: (#registerServiceProvider: using:)
 withArguments: #(scContainer aConfiguration)
"DL1-01"
"The initialization that is performed during instance creation results in the service provider simulator being available to provide service and in the service provider simulator being registered with each service category that has a matching service provider category in its service provider categories component."
sequential
 serviceRequest := nil
[] serviceProviderCategory := aConfiguration assignSPCategory
[] categoriesServed := OrderedCollection new
[] self registerServiceProvider: scContainer using:
 aConfiguration
end-sequential

category accessing

message pattern serviceProviderCategory
method properties
"Total correctness"
true **results-in** methodReturnValue = serviceProviderCategory
"DL1-02"

```

sequential
^serviceProviderCategory
end-sequential

message pattern serviceRequest
method properties
"Total correctness"
true results-in methodReturnValue = serviceRequest      "DL1-03"
sequential
^ serviceRequest
end-sequential

```

```

category testing
message pattern canAcceptNextSR: requestingServiceCategory
method properties
"Total correctness"
true results-in
  methodReturnValue = ((requestingServiceCategory =
    categoriesServed at: (categoryIndex + 1)) ^
    (serviceRequest isNil))      "DL1-04"
sequential
^ (serviceRequest isNil)
  if requestingServiceCategory =
    categoriesServed at: categoryIndex + 1 ~
^false
  if requestingServiceCategory ~=
    categoriesServed at: categoryIndex + 1
end-sequential

```

```

category modifying
message pattern registerServiceProvider: scContainer
  using: aConfiguration
"Registers the ServiceProviderSimulator with the relevant
service categories"
method macros
maxCategories ≡ aConfiguration maximumServiceCategories
method properties
"Total correctness"
true results-in methodReturnValue = self ^
  <VaServiceCategory where
    scContainer includes: aServiceCategory ^
    aServiceCategory servicedBy: serviceProviderCategory ::
    aServiceCategory postconditions: (#addSP:)
    withArguments: #(self) ^
    categoriesServed includes:
      (aServiceCategory serviceCategory)
  > ^
  nrOfCategoriesServed = categoriesServed size ^
  categoryIndex ≥ 0      "DL1-05"
sequential
nrOfCategoriesServed := 0
[] < [] j where 1 ≤ j ≤ maxCategories ::
(scContainer at: j) addSP: self \+
nrOfCategoriesServed := nrOfCategoriesServed + 1 \+
categoriesServed addLast: ((scContainer at: j) serviceCategory)
  if (scContainer at: j) servicedBy: serviceProviderCategory
>
[] categoryIndex := 0
end-sequential

```

```

message pattern processServiceRequest: aServiceRequest
method properties
"A new simulation is required each time when a new service
request is processed."
"Total correctness"
<∀ x where 0 ≤ x ∧ x < nrOfCategoriesServed ::
categoryIndex = x ∧
aServiceRequest notNil ∧
self canAcceptNextSR:(aServiceRequest serviceRequestCategory)
results-in
    methodReturnValue = self ∧
    serviceRequest = aServiceRequest ∧
    newEventRequired ∧
    categoryIndex = (x + 1) \\ nrOfCategoriesServed
>
sequential
    newEventRequired := true
[] serviceRequest := aServiceRequest
[] categoryIndex := (categoryIndex + 1) \\ nrOfCategoriesServed
end-sequential

```

"DL1-06"

category cyclic

```

message pattern p_generateEvent
method properties
"The newEventRequired attribute is not updated here. It is only
set to true once a new service request has been received."

```

```

generatingEvent ∧ serviceRequest notNil ensures
    (serviceRequest connection)
    postconditions: (#terminate:)
    withArguments: #('completed') ∧
    serviceRequest isNil ∧ ¬generatingEvent

```

"DP1-01"

"If a service provider simulator has to generate an event, it ensures that the connection currently associated with the service request is terminated and that the service provider simulator becomes available to service a new service request."

```

parallel
(serviceRequest connection) terminate: 'completed' \+
serviceRequest := nil \+
generatingEvent := false
    if generatingEvent
end-parallel

```

```

message pattern p_updateCategoryIndex: scContainer
method properties

```

```

<∀ x where 0 ≤ x ∧ x < nrOfCategoriesServed ::
categoryIndex = x ∧
(scContainer detect:[:each |
(each serviceQCategory =
(categoriesServed at: categoryIndex))
and: [each serviceQ isEmpty]] ifNone: [nil]) notNil
ensures
categoryIndex = (x + 1) \\ nrOfCategoriesServed

```

"DP1-02"

"If the serviceQ of the service category matching the categoriesServed entry at the current categoryIndex+1 is empty, the categoryIndex is incremented modulo nrOfCategoriesServed.."



```
parallel  
categoryIndex := (categoryIndex + 1) \\ nrOfCategoriesServed  
  if (scContainer detect:[:each |  
    (each serviceQCategory =  
      (categoriesServed at: (categoryIndex+1)))  
    and: [each serviceQ isEmpty]] ifNone: [nil]) notNil  
end-parallel
```

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